INF564 Compilation Mini Python

version 1 — January 26, 2024

The goal is to build a compiler for a tiny fragment of the Python language, called Mini Python in the following, to x86-64 assembly. This fragment contains Booleans, integers, strings, and lists. It is 99% compatible with Python 3. This means that Python documentation and a Python interpreter can be used as a reference when needed.

The syntax of Mini Python is described in Sec. 1. A parser is provided (for both OCaml and Java). You have to implement static type checking (Sec. 2) and code generation (Sec. 4).

1 Syntax

We use the following notations in grammars:

$\langle rule \rangle^*$	repeats $\langle rule \rangle$ an arbitrary number of times (including zero)
$\langle rule \rangle_t^\star$	repeats $\langle rule \rangle$ an arbitrary number of times (including zero), with sep-
	arator t
$\langle rule \rangle^+$	repeats $\langle rule \rangle$ at least once
$\langle rule \rangle_t^+$	repeats $\langle rule \rangle$ at least once, with separator t
$\langle rule \rangle$?	use $\langle rule \rangle$ optionally
$(\langle rule \rangle)$	grouping

Be careful not to confuse " \star " and "+" with " \star " and "+" that are Python symbols. Similarly, do not confuse grammar parentheses with terminal symbols (and).

1.1 Lexical Conventions

Comments start with # and extend to the end of line. Identifiers follow the regular expression $\langle ident \rangle$:

The following identifiers are keywords:

and	def	else	for	if	True	False
in	not	or	print	return	None	

Integer literals follow the regular expression $\langle integer \rangle$:

 $\langle integer \rangle$::= 0 | 1–9 $\langle digit \rangle^*$

String literals are written between quotes ("). There are two escape sequences: " for the character " and \n for a new line character. We assume that string literals do not contain any character $\$ beyond these two escape sequences.

In the Python language, block structure is defined by line indentation, *i.e.*, the number of spaces at the beginning of a line (it is assumed that files do not contain tabulation characters). Most of the work is done by the lexical analyzer, which produces NEWLINE, BEGIN, and END tokens corresponding to the end of lines and the increase or decrease in indentation respectively. The algorithm is as follows. The lexical analyzer maintains a stack of integers, representing successive current indentations. This stack is sorted, with the largest value at the top. Initially, the stack contains a single value, namely 0. When the lexical analyzer encounters a carriage return, it produces a NEWLINE token, then measures the indentation at the beginning of the next line, n, and compares it with the indentation at the top of the stack, m. There are three cases:

- if n = m, we do nothing;
- if n > m, we push n on the stack and we emit a second token BEGIN ;
- if n < m, then we pop until we find the value n, emitting a token END for each value strictly greater than n popped from the stack (the value n, if any, stays on top of the stack); if n does not appear on the stack, we fail with an indentation error.

In this process, empty lines, that is lines only containing spaces or comments, are ignored.

1.2 Syntax

The grammar of source files is given in Fig. 1. The entry point is $\langle file \rangle$. Associativity and priorities are given below, from lowest to strongest priority.

operation	associativity	priority
or	left	lowest
and	left	
not		
< <= > >= == i=		\downarrow
+ -	left	
* // %	left	
- (unary)		
[strongest

Note that, contrary to Python, an expression such as x < y < z is not part of our syntax.

$\langle file \rangle$::=	NEWLINE? $\langle def angle^{\star} \langle stmt angle^+$ EOF
$\langle def \rangle$::=	def $\langle ident \rangle$ ($\langle ident \rangle^{\star}$) : $\langle suite \rangle$
$\langle suite \rangle$		$\langle simple_stmt \rangle$ NEWLINE
		NEWLINE BEGIN $\langle stmt angle^+$ END
$\langle simple_stmt \rangle$		
		$\langle ident \rangle = \langle expr \rangle$
	ĺ	$\langle expr \rangle [\langle expr \rangle] = \langle expr \rangle$
		print ($\langle expr \rangle$)
		$\langle expr \rangle$
$\langle stmt \rangle$		$\langle simple_stmt \rangle$ NEWLINE
		if $\langle expr \rangle$: $\langle suite \rangle$
		if $\langle expr \rangle$: $\langle suite \rangle$ else : $\langle suite \rangle$
		for $\langle ident \rangle$ in $\langle expr \rangle$: $\langle suite \rangle$
$\langle expr \rangle$::=	$\langle const \rangle$
		$\langle ident \rangle$
		$\langle expr \rangle$ [$\langle expr \rangle$]
		$-\langle expr \rangle$
	ĺ	not $\langle expr angle$
		$\langle expr \rangle \langle binop \rangle \langle expr \rangle$
		$(ident)$ ($(expr)^*$,)
		$[\langle expr \rangle^*,]$
	ĺ	$(\langle expr \rangle)$
$\langle binop \rangle$::=	+ - * // % <= >= > < != ==
. – .		and or
$\langle const \rangle$::=	$\langle integer \rangle ~ ~ \langle string angle ~ ~ {\tt True} ~ ~ {\tt False} ~ ~ {\tt None}$

Figure 1: Grammar of Mini Python.

2 Static Typing

Though Python is a dynamically-typed language, Mini Python is simple enough to allow us some checks at compile time. These checks are the following:

1. Any function used in an expression must be

- either a function that was previously defined;
- either the function we are currently defining (a recursive function);
- one of the three built-in functions len, list, and range. (In Mini Python, print has built-in syntax and thus is not considered as a function.)

In particular, there are no mutually recursive functions in Mini Python.

- 2. The names of the functions declared with def are pairwise distinct, and distinct from len, list, and range.
- 3. Function arity must be obeyed, *i.e.*, a function defined with *n* formal parameters must be called with exactly *n* actual parameters. Functions len, list, and range all have one parameter.
- Built-in functions list and range are exclusively used in the compound expression list(range(e)).
- 5. The formal parameters of a function must be pairwise distinct.
- 6. The scope of variables is statically defined. A variable is either local to a function or global. A local variable x is introduced either as a function parameter or via an assignment x = e anywhere in the function body. The scope of a local variable extends to the full body of the function. A global variable is introduced via an assignment in the toplevel code (the code outside of function definitions at the end of the program). Hint: It is convenient to see the toplevel code as the body of a main function. This way, global variables are simply variables that are local to function main.

Note that it is not possible to shadow a variable in Mini Python. We do not (and could not) check at compile time that a variable is defined before being used.

3 Semantics

Any value in Mini Python has a *dynamic type*. This type is assigned to the value at creation time and it cannot be modified thereafter. There are five possible types: none, bool, int, string, and list. The semantics of an operation can vary according to the type of its actual parameters. In some cases, it can lead to a runtime failure. In this case, the code produced will display a message of the form

error: some message

and will terminate with exit code 1. The message does not matter but the exit code is important, as it will be used in automated tests.

Values. The none type contains a single value, noted as None. In particular, this is the value returned by a function that reaches its return point without encountering a return statement. The bool type is the Boolean type, with two possible values are False and True. The int type is for signed 64-bit integers. Type string is for character strings. Finally, type list is for lists, which are heterogeneous and possibly empty. The elements of a list can be modified in place, with the statement $x[e_1] = e_2$, but the length of a list cannot be modified.

Built-in Function print. The print statement displays the value passed as a parameter, followed by a newline. The display format is as follows:

type	affichage
none	None
bool	False, True
int	in decimal, <i>e.g.</i> 42
list	$[e_1, \ldots, e_n], e.g. [1, 2, 3]$
string	without quotes, $e.g.$ abc
	\n is printed as a newline
	" is printed as "

Boolean Condition. The statement if and the operators and, or, and not all accept operands of any type, with the following semantics:

- None, False, the integer 0, the empty string, and the empty list are interpreted as false;
- any other value is interpreted as true.

Operators and and or are lazily-evaluated: the second operand is evaluated only if needed, *i.e.* if the value of the first operand does not determine the final value. The result of e_1 and e_2 (resp. e_1 or e_2) is either the value of e_1 or the value of e_2 . For instance, 0 and [1,2] evaluates to 0 and 1 and [1,2] evaluates to [1,2].

Iteration. The statement for x in e: s first evaluates expression e, which must be of type list. Then, for each value v in this list, in order, it assigns v to the variable x and executes the statement s.

Operators. Subtraction (-), negation (unary -), multiplication (*), division (//), and modulo (%) are only defined on type int, with signed 64-bits machine arithmetic¹. Operator + is overloaded, with two parameters of the same type and the following semantics:

type of parameters	semantics
int	arithmetic addition
string	concatenation (in a new string)
list	concatenation (in a new list)
otherwise	failure

¹Python's division is actually not machine's division, but we accept this difference.

Comparison Operators. The six comparison operators (<, <=, >, >=, ==, !=) always return a Boolean value. For any comparison operator, Boolean operands are interpreted as integer values (False being 0 and True being 1). Operators == and != are defined for operands of any type, including two operands of different types. Operators <, <=, >, and >= are limited to the following cases:

type of operands	semantics
bool or int	arithmetic comparison
string	lexicographic comparison
list	lexicographic comparison

The comparison is structural: when comparing lists, the comparison operation is recursively applied to the list elements.

Built-in Functions. Function len is only defined on types string and list. It returns the length of the string and of the list, respectively. The expression list(range(e)) is defined for an expression e that evaluates to some integer $n \ge 0$. It returns a list of integers $[0, 1, \ldots, n-1]$.

Differences w.r.t Python. There are some small runtime differences between Python and Mini Python, including (but not exhaustively):

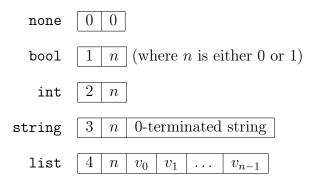
	Python	Mini Python
arithmetic	unbounded	signed 64 bits
string display within lists	['abc']	[abc]
multiplying a string/list and an integer	defined	undefined
access $l[i]$ with $i < 0$	defined	undefined

Your compiler will not be tested on programs for which there is a distinct runtime behavior between Python and Mini Python.

4 Code Generation

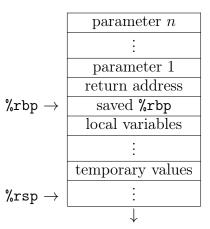
The aim is to produce a simple but correct compiler. In particular, we do not attempt to do any kind of register allocation, but simply use the stack to store any intermediate calculations. Of course, it is possible, and even desirable, to use some x86-64 registers locally. Memory is allocated using malloc and no attempt will be made to free memory.

Value Representation. Though Python is an object-oriented language, we are not compiling Mini Python using objects. We propose a simple compilation scheme (but you are free to use any other) where a value is always a pointer to some heap-allocated block. The first word (64 bits) of this block is an integer tag that encodes the type.



where n is a 64-bit integer and v_i are values (*i.e.*, pointers to other heap-allocated blocks). For strings and lists, the integer n is the length. For a string, the third component is a 0-terminated string stored in n + 1 bytes (we assume ASCII strings in Mini Python).

Stack Layout. We suggest a compilation scheme where all parameters are passed on the stack (each of them being a 64-bit pointer), and where the return value is in register **%rax**. The stack frame is as follows:



Local variables are allocated on the stack. The top of the stack is used to store intermediate computations, such as the value of e_1 during the evaluation of e_2 in a binary operation $e_1 \oplus e_2$.

With recent versions of the libc, it is important to have a 16-byte stack alignment when calling library functions such as malloc or printf (this is required by the System V Application Binary Interface). Since it is not always easy to ensure stack alignment when calling library functions (because of intermediate computations temporarily stored on the stack), it may be convenient to introduce wrappers around library functions, as follows:

```
my_malloc:
```

```
pushq %rbp
movq %rsp, %rbp
andq $-16, %rsp # 16-byte stack alignment
call malloc
movq %rbp, %rsp
popq %rbp
ret
```

These wrappers are simply concatenated to the generated assembly code — and of course any call to malloc is replaced with a call to my_malloc.

Here is a list of functions from the C standard library that you may want to use (feel free to use any other):

```
void *malloc(size_t size);
    malloc(n) returns a pointer to a freshly heap-allocated block of size n
    You don't have to free memory.
int putchar(int c);
    putchar(c) writes the character c to standard output (ignore the return value)
int printf(const char *format, ...);
    printf(f,...) write to standard output according to the format string
    (ignore the return value). Register %rax must be set to zero before calling printf.
int strcmp(const char *s1, const char *s2);
    compare strings s1 and s2, returning 0 if they are equal, a negative value
    if s1 is smaller than s2, and a positive value if s1 is greater
char *strcpy(char *dest, const char *src);
    copy the 0-terminated string src to dest, including the '0' character
    (ignore the return value)
char *strcat(char *dest, const char *src);
    appends the 0-terminated string src at the end of string dest, assuming there is
    enough space (ignore the return value)
```

Important Notice. Grading involves (for one part only) some automated tests using small Python programs with print commands. They are compiled with your compiler, and the output is compared to the expected output. This means you should be careful in compiling calls to print.

5 Project Assignment (due March 17, 6pm)

The project must be done alone or in pair, in Java or OCaml. It must be delivered on Moodle, as a compressed archive containing a directory with your name(s) (*e.g.* dupont-durand). Inside this directory, source files of the compiler must be provided (no need to include compiled files). The command make must create the compiler, named minipython. The compilation may involve any tool (such as dune for OCaml) and the Makefile can be as simple as a call to such a tool. The command minipython may be a script to run the compiler, for instance if the compiler is implemented in Java.

The archive must also contain **a short report** explaining the technical choices and, if any, the issues with the project and the list of whatever is not delivered. The report can be in format ASCII, Markdown, or PDF.

The command line of minipython accepts an option (among --parse-only and --type-only) and exactly one file with extension .py. If the file is parsed successfully, the compiler must terminate with code 0 if option --parse-only is on the command line. Otherwise, the compiler moves to static type checking. Any type error must be reported as follows:

file.py:4:6: bad arity for function f

The location indicates the filename name, the line number, and the column number. Feel free to design your own error messages. The exit code must be 1.

If the file is type-checked successfully, the compiler must exit with code 0 if option --type-only is on the command line. Otherwise, the compiler generates x86-64 assembly code in file file.s (same name as the input file, but with extension .s instead of extension .py). The x86-64 file will be compiled and run as follows

gcc file.s -o file
./file

possibly with option -no-pie on the gcc command line. Any runtime error must be reported, but no location nor a detailed message is expected so it is fine to simply output

error

and terminate with exit code 1.

An Extension of Your Choice. Last, but not least, you have to implement an extension of your choice. This can be

- the support of another Python construct;
- more static analysis;
- awesome error messages;
- a compiler optimization;
- etc.

This extension must be described in the report and illustrated with test files.